

Efficiently computing the minimum rank of a matrix in a monoid of zero-one matrices

joint work with Stefan Kiefer

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Zero-one matrix monoids

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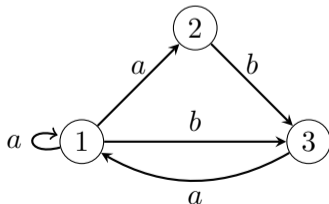
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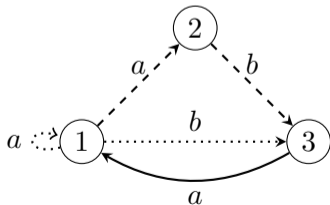
$$M(a) = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \end{pmatrix} \quad M(b) = \begin{pmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$



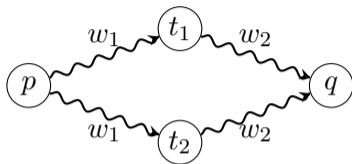
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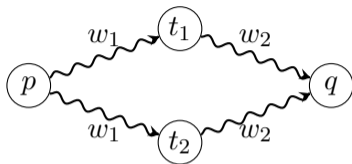


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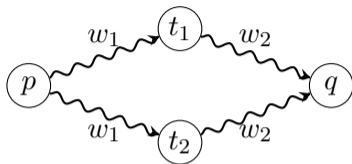
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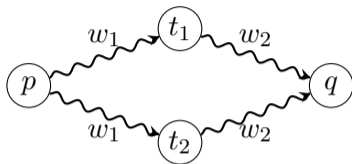
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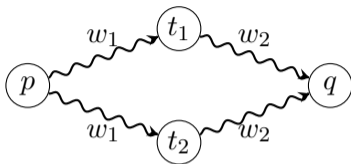
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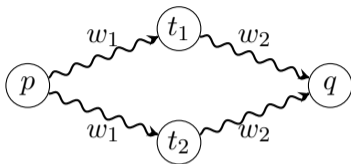
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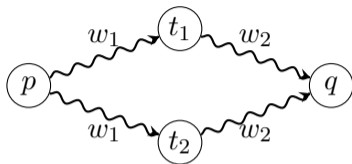
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- Indeed, it's enough to check if $\bar{A}x = x$ has a non-zero solution.

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- Similarly, the rank conjecture generalises it to products of length $(n - r)^2$ in the case of minimum rank $r \geq 1$.
- Minimum rank of a matrix is computable in polynomial time (Carpi, TCS 1988 + Kiefer, Mascle, STACS 2019).

Why are these problems nice?



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- It's also “combinatorial”, while Kiefer-Masole treatment of the zero-rank case is linear-algebraic.
- If the minimum rank is positive, can we still compute it by linear algebra (in NC^2)?— Wasn't known even for DFAs.

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Minimum rank in NC^2

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Given a zero-one matrix monoid, one can compute the minimum rank of a matrix in it in NC^2 .

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- By using this result, we define nicely behaving weights of states in the unambiguous automaton.
- Some sort of Kraft–McMillan (in)equality?

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- Can one find a matrix of minimum rank in a zero-one matrix monoid in NC? Not known even for complete DFAs (one 1 per row).
- Given two matrices generating a zero-one matrix monoid, can we multiply them faster than in $\mathcal{O}(n^\omega)$?
- If \mathcal{M} contains the zero matrix, Kiefer and Mascle showed that this matrix can be represented as a product of at most n^5 generators. The best known lower bound is $n(n+1)/2$. Can we do anything about that?

Thank you!