Local equivalence of stabilizer states

a graphical characterisation

Nathan Claudet and Simon Perdrix

STACS 2025 - 06/03/25 arXiv:2409.20183



Graph states, local unitary equivalence, local Clifford equivalence & local complementation



¹Edges do not have a direction. ²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.



¹Edges do not have a direction.

²No multiples edges and no loops.

A graph state is a quantum state represented by an undirected¹ and simple² graph. The vertices represent the qubits and the edges represent entanglement.



Graph states \sim Stabilizer states.

¹Edges do not have a direction.

²No multiples edges and no loops.

Graph states are useful entangled resources (measurement-based quantum computation, quantum error correction...). \rightarrow It is a fundamental problem to know whether two graph states have the same entanglement.

Entanglement of graph states

Graph states are useful entangled resources (measurement-based quantum computation, quantum error correction...). \rightarrow It is a fundamental problem to know whether two graph states have the same entanglement. For graph states, having the same entanglement = being **local unitary** equivalent, i.e. related by single-qubit unitary quantum gates.



Entanglement of graph states

Graph states are useful entangled resources (measurement-based quantum computation, quantum error correction...). \rightarrow It is a fundamental problem to know whether two graph states have the same entanglement. For graph states, having the same entanglement = being **local unitary** equivalent, i.e. related by single-qubit unitary quantum gates.



An easier subproblem: local Clifford equivalence

Two graph states are said **local Clifford equivalent** (or LC-equivalent) if they are related by unitaries in the local Clifford group.



An easier subproblem: local Clifford equivalence

Two graph states are said **local Clifford equivalent** (or LC-equivalent) if they are related by unitaries in the local Clifford group.



Theorem (Van den Nest, Dehaene, De Moor, 2004)

Two graph states are local Clifford equivalent iff the two corresponding graphs are related by **local complementations**.

Local complementation

Definition (Kotzig, 1966)

A local complementation on a vertex u consists in complementing the (open) neighbourhood of u.



Local complementation

Definition (Kotzig, 1966)

A local complementation on a vertex u consists in complementing the (open) neighbourhood of u.



Local complementation

Definition (Kotzig, 1966)

A local complementation on a vertex u consists in complementing the (open) neighbourhood of u.



There exists an efficient algorithm (Bouchet, 1991) to recognise whether two graphs are related by local complementations, implying an efficient algorithm to recognise whether two graph states are local Clifford equivalent.

$\mathsf{LU} \neq \mathsf{LC}$

Unfortunately, LU \neq LC, i.e. local Clifford equivalence and local unitary equivalence do not coincide.

$\mathsf{LU} \neq \mathsf{LC}$

Unfortunately, LU \neq LC, i.e. local Clifford equivalence and local unitary equivalence do not coincide. \leftarrow 27-qubit pair of graph states that are local unitary equivalent but not local Clifford equivalent (Ji et al. 2008).



$\mathsf{LU} \neq \mathsf{LC}$

Unfortunately, LU \neq LC, i.e. local Clifford equivalence and local unitary equivalence do not coincide. \leftarrow 27-qubit pair of graph states that are local unitary equivalent but not local Clifford equivalent (Ji et al. 2008).



Consequence: local complementation does **not** capture the local unitary equivalence of graph states.

• Graph states over at most 8 qubits (Cabello et al. 2009)

- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs (Van den Nest, Dehaene, De Moor, 2005)

- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs (Van den Nest, Dehaene, De Moor, 2005)
- Complete bipartite graphs (Tzitrin, 2018)

- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs (Van den Nest, Dehaene, De Moor, 2005)
- Complete bipartite graphs (Tzitrin, 2018)
- Graphs with no cycle of length 3 or 4 (Zeng et al. 2007)

- Graph states over at most 8 qubits (Cabello et al. 2009)
- Complete graphs (Van den Nest, Dehaene, De Moor, 2005)
- Complete bipartite graphs (Tzitrin, 2018)
- Graphs with no cycle of length 3 or 4 (Zeng et al. 2007)

But what about local unitary equivalence for **any** graph ? Can we construct a graphical characterisation ?

Generalising local complementation to capture local unitary equivalence













A sequence of local complementations may leave the graph invariant.



A **2-local complementation** consists in toggling every edge that was toggled 2 mod 4 times by the idempotent local complementations. (There are also some additional conditions on the edges for the 2-local complementation the be valid.)

A sequence of local complementations may leave the graph invariant.



A **2-local complementation** consists in toggling every edge that was toggled 2 mod 4 times by the idempotent local complementations. (There are also some additional conditions on the edges for the 2-local complementation the be valid.)

3-local complementation is a refinement of idempotent 2-local complementation, and so on...

 \rightarrow Infinite family of graphical operations parametrised by an integer r:

r-local complementations

1-local complementation = local complementation.

Theorem (this work)

Two graph states are local unitary equivalent iff the two corresponding graphs are related by r-local complementations for some r.

An infinite hierarchy of local equivalences



An infinite hierarchy of local equivalences



An infinite hierarchy of local equivalences



Proof sketch: Minimal local set



Minimal local sets are subsets of vertices that are invariant by local unitary equivalence and carry information on the possible local unitaries that maps graph states to other graph states.

Theorem (<u>C</u>, Perdrix, 2024)

Each vertex of a graph is covered by at least one minimal local set.

Proof sketch: Standard form for graph states



Proof sketch: Standard form for graph states



Proof sketch: Standard form for graph states



Applications

LU=LC for some repeater graph states, as conjectured in [Tzitrin, 2018].



Applications

LU=LC for some repeater graph states, as conjectured in [Tzitrin, 2018].



Theorem (<u>C</u>, Perdrix, 2025) LU=LC for graph states up to 19 qubits.

It was previously known that LU=LC for graph states up to 8 qubits, and there exists a 27-qubit pair for which LU \neq LC.

Applications

LU=LC for some repeater graph states, as conjectured in [Tzitrin, 2018].



Theorem (<u>C</u>, Perdrix, 2025)

LU=LC for graph states up to 19 qubits.

It was previously known that LU=LC for graph states up to 8 qubits, and there exists a 27-qubit pair for which LU \neq LC.

Theorem (C, Perdrix, 2025)

There exists an algorithm that decides whether two graph states are local unitary equivalent with runtime $n^{\log_2(n)+O(1)}$.

Summary

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

Open questions:

Does there exist a counter-example to LU=LC between 20 and 26 qubits ?

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

Open questions:

- Does there exist a counter-example to LU=LC between 20 and 26 qubits ?
- Does there exist a polynomial-time algorithm for local unitary equivalence ?

Thanks



arXiv:2409.20183